

# CS 4510 : Automata and Complexity

## Simulating NTMs using DTMs

Subrahmanyam Kalyanasundaram

February 16, 2010

### Simulation of NTMs using DTMs

**Theorem 1.** *Every nondeterministic Turing machine has an equivalent deterministic Turing machine.*

We need to simulate a nondeterministic Turing machine (NTM) by a deterministic Turing machine (DTM). Given an NTM  $N$ , we need to simulate it using a DTM  $D$  on an input  $w$ .  $N$  is a 1-tape NTM. We shall use a 3-tape DTM  $M$  to simulate it. Since we have seen that multitape DTMs are equivalent to single tape DTMs, this is good enough.

We look at the configuration graph of  $N$  to understand the simulation. By the definition of the NTM, there are a finite set of choices that  $N$  can move from each given configuration. Call this number  $b$ . Order the possible successor configurations of a given configuration with  $1, 2, \dots, b$ . This gives us the configuration tree of  $N$ , each node having at most  $b$  children. The DTM  $D$  simulates  $N$  by doing a breadth first search on the configuration graph. The tapes of  $D$  contain the following information.

1. Input  $w$
2. The simulation tape. This contains the contents of  $N$ 's tape as it is working.
3. The counter for the node in the tree.

Please see the next page for the algorithm described formally.

### Guess and verify NTMs

If we have an NTM which is a decider, i.e., one which is known to halt within  $t$  steps, we could make all the guesses (nondeterministic choices) right at the beginning of the computation. We could guess one of the strings  $p$  from  $\{1, 2, \dots, b\}^t$  and proceed the computation guided by the string  $p$ . At each step, we follow the respective symbol of  $p$ , when faced with choices. Notice that once we make the guess of  $p$ , the remaining computation is completely **deterministic**. The machine has to just look up the corresponding symbol in  $p$  and decide on the choice to move. The machine accepts if the computation leads to an accept, and rejects if it leads to a reject. The configuration graph of this “guess and verify” NTM, it would look like a tree with a paths attached to each of the leaves.

---

**Algorithm 1** Simulating NTM using DTM. Given NTM  $N$  and input string  $w$ .

---

```
1: Tape 3  $\leftarrow \varepsilon$ 
2:  $l = 0$  //  $l$  corresponds to the length of the counter in tape 3
3:  $l \leftarrow l + 1$ ; Tape 3  $\leftarrow 1^l$ 
4: Set Rejectedsofar to TRUE
5: Tape 2  $\leftarrow$  Tape 1
6: Simulate  $N$  on contents of tape 2, guided by contents of tape 3.
   Use tape 3 contents to choose from the options available and go to the required node
   • Current configuration is not defined or rejecting: GOTO step 7
   • Current configuration is accepting: HALT and ACCEPT
   • Current configuration is defined but neither accepting nor rejecting:
     Rejectedsofar = FALSE; GOTO step 7
7:
8: if Tape 3 =  $b^l$  then
9:   if Rejectedsofar = TRUE then
10:    HALT and REJECT
11:   else
12:    GOTO step 3
13: else
14:   Increment string in Tape 3 and GOTO step 5
```

---